

# FURTHER TRIANGLE TILINGS

CHAIM GOODMAN-STRAUSS

# Further Triangle tilings

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## Abstract

Here we give a more complete reckoning of the conjecture discussed in “Regular Production Systems and Triangle Tilings” [16]

Here we discuss which triangles do, and which don’t, admit a tiling of  $\mathbb{H}^2$ ,  $\mathbb{E}^2$ , and  $\mathbb{S}^2$ . These notes are meant to pick up as “Regular Production Systems and Triangle Tilings” [16] leaves off.

## 1 A little notation for cases

To simplify our analysis and conflate cases we adopt the following notation. A “form”  $|F|$  denotes a family of sets of planes in  $\Pi$ ; these families will be invariant under the symmetry of  $\mathcal{O}$ . Unfortunately, it will be possible to denote a given family in many ways.

A form will consist of rows of three symbols each. The symbols are  $*$ ,  $d$ ,  $e$ ,  $0/1$ , standing respectively for any natural number, any odd natural, any even natural, and either 0 or 1; and for any fixed natural numbers  $n, m$  we also have  $n, n^+, n^{+m}$  standing for  $n$ , any  $n + k$ , any  $n + 2k$ , any  $n + mk$ , respectively, with  $k \geq 0$ . In addition a row may be demarked by  $\natural$ .

A plane  $\pi_{rst}$  is of the form of a given row  $xyz$  if and only if  $x, y, z$  can stand for  $r, s, t$ . A set  $P \subset \Pi$  is an element of a form  $F$  if and only if: There exists a  $Q$  (possibly  $P$  itself) such that (a)  $Q = \rho P$  for some permutation  $\rho$  of the coordinates of  $\mathcal{O}$ ; (b) every plane in  $Q$  is of the form of some row of  $F$ ; and (c) every non-demarked row is the form of exactly one plane in  $Q$  (demarked forms may be represented once, or more, or not at all, in  $Q$ ).

So for example, letting  $P = \{\pi_{241}, \pi_{357}\}$ ,  $P \in \left| \begin{smallmatrix} e & 2^+ & 1 \\ 3^+ & 3^+ & 1+3 \end{smallmatrix} \right|$ , among many other forms.

Note we can denote a family of sets of planes in many ways; for example:

$$\left| \begin{smallmatrix} 2 & 6 & 3 \\ 0 & 8 & 6 \end{smallmatrix} \right| = \left| \begin{smallmatrix} 8 & 6 & 0 \\ 6 & 3 & 2 \end{smallmatrix} \right|.$$

Given a form  $|F|$ , we will write

$$[F] := \text{“for all } P \in |F|, [P]\text{”}$$

$$(F) := \text{“for all } P \in |F|, (P)\text{”}$$

We will also take  $\mathcal{V}(F) := \{\mathcal{V}(P)\}$ . Recall that the notation was chosen so that the statements  $[P], (P)$  refer to tilings by  $T \in P$ ; consequently, it is perfectly possible that  $(F)$  does not hold even though there is no 2-nion, say, by any  $V \in \mathcal{V}(F)$ . This device allows us to distinguish between the statements “there is a tiling using only such-and-such vertex arrangements” and “there is a tiling by triangles whose angles satisfy such-and-such system of equations”. Lemma 1.1 helps a great deal as we analyze cases later.

**Lemma 1.1** 1.  $(\natural 0/1 * *)$ , with the exception of the scalene tetrahedra  $[1 1 1]$ .

2.  $\left( \begin{smallmatrix} \natural & 2 & 0 & 2^+ \\ \natural & 0/1 & * & * \end{smallmatrix} \right)$  with the exceptions of  $\left[ \begin{smallmatrix} 2 & 0 & 2 \\ 0 & 2^+ & 0 \end{smallmatrix} \right], \left[ \begin{smallmatrix} 2 & 0 & 2 \\ 1 & 3^+ & 1 \end{smallmatrix} \right]$ .

3.  $\left( \begin{smallmatrix} 2 & 0 & 0 \\ \natural * & * & * \end{smallmatrix} \right)$  with the exception that for all non-negative  $p$ , positive  $q$ :

$$\left[ \begin{smallmatrix} 2 & 0 & 0 \\ 0 & 2p & 2q \\ 0 & 2q & 2p \end{smallmatrix} \right] \text{ and } \left[ \begin{smallmatrix} 2 & 0 & 0 \\ 1 & 2p+1 & 2q+1 \\ 1 & 2q+1 & 2p+1 \end{smallmatrix} \right].$$

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We'll first construct one  $\mathcal{V}'$ -complex. Replace every edge in  $\mathcal{E}$  with a lune as shown below:



The orientation of the lune does not matter particularly. The three labels  $x, y, z$  are  $A, B, C$  (or their reflections, depending on the orientation), but we are free to assign which is which so long as the labels on the outside edges are compatible with the edge in  $\mathcal{E}$ .

Now simply take the universal cover of the two-fold branched cover through all the marked points in the interior of all the lunes. The resulting complex will be a  $\mathcal{V}'$ -complex. Each lune will become a strip as shown at right above.

But in fact, there are far more choices available when  $|\mathcal{E}| \geq 2$ . Let  $\mathcal{C}_1$  be the complex formed from  $\mathcal{C}$  by doubling each edge of  $\mathcal{E}$ , creating a lune shaped hole bounded by two copies of the original edge. Let  $\mathcal{C}_2 = \tilde{\mathcal{C}}_1$  be the universal cover of  $\mathcal{C}_1$ .  $\mathcal{C}_2$  now has either two, or infinitely many, boundary components. Each component may be marked  $A, B, C$  depending on the markings of the original edge in  $\mathcal{E}$ . Define three strips  $s_A, s_B, s_C$  of triangles:  $s_x$  is shown above.

We will not make this next bit completely precise; to do so only requires a careful induction. Begin with a copy of  $\mathcal{C}_2$  and to each boundary component labeled  $X = A, B, C$ , attach a copy of  $s_X$ , making an arbitrary choice of orientation. Next, to each boundary component, adjoining some  $s_X, X = A, B, C$  of this new complex attach a copy of  $\mathcal{C}_2$ , along a boundary labeled  $X$ , again with an arbitrary choice of orientation. At each stage, repeat this process, ad infinitum.

Now note that every vertex in the final complex is of the form of one of  $\mathcal{V}'$  and we have a  $\mathcal{V}'$ -complex.  $\square$

We give a quick application or two:

Note that in this way,  $\begin{bmatrix} 2 & 0 & 0 \\ 1 & 2p+1 & 2q+1 \\ 1 & 2q+1 & 2p+1 \end{bmatrix}$  can be produced from  $\begin{bmatrix} 2 & 0 & 0 \\ 0 & 2p & 2q \\ 0 & 2q & 2p \end{bmatrix}$ .

Similarly, by choosing appropriate edges in  $\begin{bmatrix} 2 & 0 & 2 \\ 0 & 2+ & 0 \end{bmatrix}$  we have:

**Corollary 2.2**  $\begin{bmatrix} 3 & 3 & 1 \\ 0 & 0 & 2+ \end{bmatrix}$  and  $\begin{bmatrix} 3 & 3 & 1 \\ 1 & 1 & 3+ \end{bmatrix}$

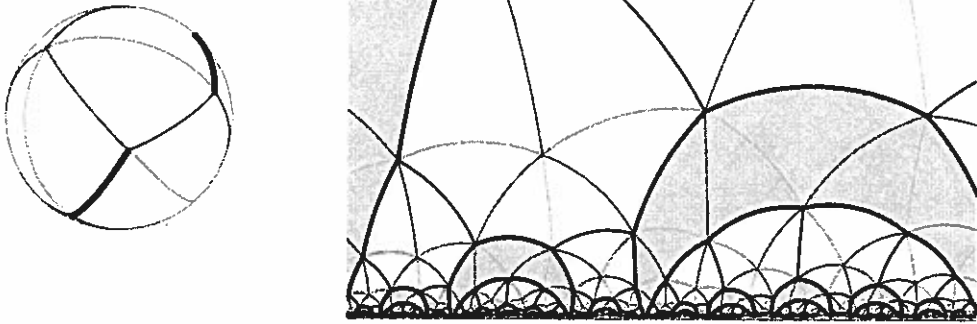


Figure 2: A tiling by  $\mathcal{V} \begin{pmatrix} 6 & 0 & 0 \\ 1 & 3 & 3 \end{pmatrix}$ ; this can be viewed as arising from  $\mathcal{V} \begin{pmatrix} 6 & 0 & 0 \\ 0 & 2 & 2 \end{pmatrix}$  by an application of Lemma 2.1. The tiling shown has symmetry 2223 in the Conway orbifold notation; however, there are uncountably many tilings by this tile, most of which have no symmetry.

We can use branched covers in other ways as well:

**Theorem 2.3** For all  $m, n$ :  $\begin{pmatrix} n & m & 0 \\ 1 & 1 & 3 \end{pmatrix}$ , with the exception that for all even  $n \geq 4$ ,  $\begin{bmatrix} n & n & 0 \\ 1 & 1 & 3 \end{bmatrix}$

**Proof** First, let  $n$  be even and positive. Then letting  $\mathcal{V} = \mathcal{V} \begin{pmatrix} n & n & 0 \\ 1 & 1 & 3 \end{pmatrix}$ , construct a  $\mathcal{V}$ -complex by taking the universal cover of the indicated branched cover of the graph in Figure 3.

### 3 Theorems

Before delving into the combinatorial thicket of the space of triangles, we pause to illustrate these techniques with a well-known theorem:

#### 3.1 The Poincaré triangle theorem

**Theorem 3.1 (Poincaré)**  $\begin{bmatrix} 4^+ & 0 & 0 \\ 0 & 4^+ & 0 \\ 0 & 0 & 4^+ \end{bmatrix}$ .

In fact we only prove here  $\begin{bmatrix} 6^+ & 0 & 0 \\ 0 & 6^+ & 0 \\ 0 & 0 & 6^+ \end{bmatrix}$ .

**Proof** Let  $P = \{\pi_{(2p)00}, \pi_{0(2q)0}, \pi_{00(2r)}\}$  and let  $T \in \cap P$ . Let  $\mathcal{V} = \mathcal{V}(P)$ . Then certainly the three vertex configurations  $(\alpha\bar{\alpha})^p, (\beta\bar{\beta})^q, (\gamma\bar{\gamma})^r$  lie in  $\mathcal{V}$  and the following rules lie in  $\mathcal{R}_{\mathcal{V}}$ :

$(\alpha\bar{\alpha})^p$	$(\beta\bar{\beta})^q$	$(\gamma\bar{\gamma})^r$
$[\overline{BB}] \mapsto (\overline{AA})^{i+1}\overline{\beta C}$	$[\overline{CC}] \mapsto (\overline{BB})^{j+1}\overline{\gamma A}$	$[\overline{AA}] \mapsto (\overline{CC})^{k+1}\overline{\alpha B}$
$[\overline{CC}] \mapsto (\overline{AA})^{i+1}\overline{\gamma B}$	$[\overline{AA}] \mapsto (\overline{BB})^{j+1}\overline{\alpha C}$	$[\overline{BB}] \mapsto (\overline{CC})^{k+1}\overline{\alpha A}$
$[\overline{B\alpha C}] \mapsto A(\overline{AA})^i\overline{\gamma B}$	$[\overline{C\beta A}] \mapsto B(\overline{BB})^j\overline{\alpha C}$	$[\overline{A\gamma B}] \mapsto C(\overline{CC})^k\overline{\beta A}$
$[\overline{C\alpha B}] \mapsto \overline{A}(\overline{AA})^i\overline{\beta C}$	$[\overline{A\beta C}] \mapsto \overline{B}(\overline{BB})^j\overline{\gamma A}$	$[\overline{B\gamma A}] \mapsto \overline{C}(\overline{CC})^k\overline{\alpha B}$

where  $i = p - 3, j = q - 3, k = r - 3$ ; recall our convention for abbreviating the words in  $\mathcal{L}$ . For example,  $(\overline{AA})^2\overline{\beta C} = \overline{AA}\overline{AA}\overline{\beta C} = [\overline{AA}][\overline{AA}][\overline{A\beta C}]$ .

Then taking  $\mathcal{A}^-$  to be the twelve letters on the left of each rule, taking  $\mathcal{L}^- = \mathcal{L}|_{\mathcal{A}^-}$  (that is,  $\mathcal{L}$  restricted to the letters of  $\mathcal{A}^-$ ) and  $\mathcal{R}^-$  to be specified by the twelve rules above, one can easily check that  $(\mathcal{A}^-, \mathcal{L}^-, \mathcal{R}^-)$  forms a cut of  $(\mathcal{A}, \mathcal{L}, \mathcal{R}_{\mathcal{V}})$ . Consequently, there is a orbit in  $\mathcal{L}^\infty$  under  $\mathcal{R}$ , there is a  $\mathcal{V}$ -complex, and finally,  $T$  admits a tiling.  $\square$

Indeed, in some ways, this is precisely the core of Poincaré's construction.

#### 3.2 A quick lemma

The following lemma is quite helpful and generalizes readily.

**Lemma 3.2** *Let  $\mathcal{V}$  be a set of vertex arrangements. Then  $(\mathcal{A}, \mathcal{L}, \mathcal{R}_{\mathcal{V}})$  has a cut if there exists a function  $f : \{A, B, C, \overline{A}, \overline{B}, \overline{C}\} \rightarrow \{A, B, C, \overline{A}, \overline{B}, \overline{C}\}$  such that there is a non-empty collection of letters  $\mathcal{A}^- \subset \mathcal{A}$  and rules  $\mathcal{R}^- \subset \mathcal{R}_{\mathcal{V}}$ , one or more for each letter in  $\mathcal{A}^-$ , each of the form  $[X \dots Y] \xrightarrow{\mathcal{R}^-} [f(X) \dots] \dots [\dots f(Y)]$  with letters in  $\mathcal{A}^-$ .*

**Proof** This is a trivial application of Lemma 2.9 of [16], taking one rule for each letter.  $\square$

#### 3.3 On 2-dimensional affine subspaces in $\cup\Pi$

Recalling that  $(dde), (dee), (0/1 \ **)$  with the exception of [111], this Proposition completely settles the question of whether or not a given triangle, lying on exactly one plane of  $\Pi$ , admits a tiling. That is, we now have given necessary and sufficient conditions for admitting a tiling, for a measure-one set of triangles in  $\cup\Pi$ .

**Proposition 3.3**  $[2^+ \ 2^+ \ 2^+]$  and  $[3^+ \ 3^+ \ 3^+]$

Please note that in throughout the remainder of the paper we will be content to find *one* cut or lift; but the particular choice we make is rather arbitrary.

**Proof** We simply demonstrate that one may take the identity function for  $f$  in an application of Lemma 3.2. In each case we'll simply list the rules in  $\mathcal{R}^-$ ; the alphabet  $\mathcal{A}^-$  will be given as the

$$\begin{array}{cccc} \left| \begin{array}{ccc} 3^+ & 3^+ & 1 \\ 3^+ & 1 & 3^+ \end{array} \right| & \left| \begin{array}{ccc} 3^+ & 3^+ & 1 \\ 1 & 1 & 3^+ \end{array} \right| & \left| \begin{array}{ccc} 3^+ & 3^+ & 1 \\ 2^+ & 0 & 2^+ \end{array} \right| & \left| \begin{array}{ccc} 3^+ & 3^+ & 1 \\ 0 & 0 & 2^+ \end{array} \right| \\ \\ \left| \begin{array}{ccc} 3^+ & 1 & 1 \\ 0 & 2^+ & 2^+ \end{array} \right| & \left| \begin{array}{ccc} 2^+ & 2^+ & 0 \\ 2^+ & 0 & 2^+ \end{array} \right| & \left| \begin{array}{ccc} 2^+ & 2^+ & 0 \\ 0 & 0 & 2^+ \end{array} \right| & \end{array}$$

We will have at least a little to say about each of these forms.

### 3.5 The form $\left| \begin{array}{ccc} 3 & 3 & 1 \\ 3 & 1 & 3 \end{array} \right|$

There are cuts for  $\left| \begin{array}{ccc} 3^+ & 5^+ & 1 \\ 3^+ & 1 & 5^+ \end{array} \right|$ . Apparently there is no cut of  $\left| \begin{array}{ccc} 3 & 3 & 1 \\ 3 & 1 & 3 \end{array} \right|$ , but there is a lift. The cases  $\left| \begin{array}{ccc} 5^+ & 3 & 1 \\ 3^+ & 1 & 3^+ \end{array} \right|$  and  $\left| \begin{array}{ccc} 3^+ & 3 & 1 \\ 3^+ & 1 & 5^+ \end{array} \right|$  have been vexing and remain open.

#### Theorem 3.4 $\left[ \begin{array}{ccc} 3 & 3 & 1 \\ 3 & 1 & 3 \end{array} \right]$

**Proof** Define new symbols  $C, \bar{C}, \bar{C}$ , mapping to  $C, \bar{C}$ , and  $\gamma, \bar{\gamma}$  mapping to  $\gamma, \bar{\gamma}$ ; define  $\mathcal{A}^+, \mathcal{L}^+$ , and  $\mathcal{R}^+$  as indicated below—the letters of  $\mathcal{A}^+$  are on the right; the language follows the same method as  $\mathcal{L}_T$  (adjacent symbols must match) and the rules are as given. It is not hard to see that these form a lift for  $\mathcal{A}_\nu, \mathcal{L}_\nu, \mathcal{R}_\nu$ .

$$\begin{array}{llll} C\bar{C} \rightarrow C\bar{C}C\gamma C & \beta\bar{\beta}\beta\bar{\gamma}\bar{\gamma}\alpha & C\bar{C} \rightarrow \bar{C}C\bar{C}\bar{\gamma}\bar{C} & \beta\bar{\beta}\bar{\alpha}\bar{\gamma}\bar{\gamma}\bar{\beta} \\ \bar{C}C \rightarrow C\bar{C}C\gamma C & \bar{\alpha}\alpha\beta\bar{\gamma}\bar{\gamma}\alpha & \bar{C}C \rightarrow \bar{C}C\bar{C}\bar{\gamma}\bar{C} & \bar{\alpha}\alpha\bar{\alpha}\bar{\gamma}\bar{\gamma}\bar{\beta} \\ C\gamma C \rightarrow C\bar{C}\bar{\gamma}\bar{C} & \beta\gamma\alpha\bar{\alpha}\bar{\gamma}\alpha & C\gamma C \rightarrow \bar{C}C\gamma C & \beta\gamma\alpha\beta\bar{\gamma}\bar{\beta} \\ \bar{C}\bar{\gamma}\bar{C} \rightarrow C\bar{C}\bar{\gamma}\bar{C} & \bar{\alpha}\bar{\gamma}\bar{\beta}\bar{\alpha}\bar{\gamma}\alpha & \bar{C}\bar{\gamma}\bar{C} \rightarrow \bar{C}C\gamma C & \bar{\alpha}\bar{\gamma}\bar{\beta}\beta\bar{\gamma}\bar{\beta} \end{array}$$

□

#### Theorem 3.5 $\left[ \begin{array}{ccc} 3^+ & 5^+ & 1 \\ 3^+ & 1 & 5^+ \end{array} \right]$

**Proof** We take the following cut for  $\left[ \begin{array}{ccc} 2p+3 & 2q+5 & 1 \\ 2r+3 & 1 & 2s+5 \end{array} \right]$ ,  $p, q, r, s \geq 0$ . Note that this is in fact a cut by applying Lemma 3.2 using the function  $f(A, B, C, \bar{C}, \bar{B}, \bar{A}) = (C, B, B, \bar{C}, \bar{C}, \bar{B})$ . The rules are:

$$\begin{array}{lll} [BA] \rightarrow B\bar{B}^\rho BA\bar{A}\bar{\alpha}C & [B\bar{B}] \rightarrow B\bar{B}^\rho B\bar{B}B\beta\bar{C} & [CB] \rightarrow B\bar{B}^\rho BA\bar{A}\beta B \\ [C\bar{C}] \rightarrow B\bar{B}^\rho \bar{C}AA\bar{\gamma}\bar{C} & [\bar{A}A] \rightarrow \bar{B}^\rho \bar{C}AA\bar{A}\bar{\alpha}C & [\bar{A}B] \rightarrow \bar{B}^\rho \bar{C}A\bar{B}B\beta\bar{C} \\ [B\beta\bar{C}] \rightarrow B\bar{B}^\rho BA\bar{\gamma}\bar{C} & [\bar{A}\beta B] \rightarrow \bar{B}^\rho BA\bar{A}\beta B & [\bar{A}\bar{\alpha}C] \rightarrow \bar{B}^\rho B\bar{B}\bar{C}\bar{\gamma}B \\ \\ [AC] \rightarrow C^\rho BAC\bar{C}\bar{\gamma}B & [A\bar{A}] \rightarrow C^\rho BA\bar{A}\bar{A}\alpha\bar{B} & [\bar{B}B] \rightarrow \bar{C}C^\rho BA\bar{A}\beta B \\ [\bar{B}C] \rightarrow \bar{C}C^\rho \bar{C}AA\bar{\gamma}\bar{C} & [\bar{C}C] \rightarrow \bar{C}C^\rho \bar{C}C\bar{C}\bar{\gamma}B & [CA] \rightarrow \bar{C}C^\rho \bar{C}AA\alpha\bar{B} \\ [A\alpha\bar{B}] \rightarrow C^\rho \bar{C}CB\beta\bar{C} & [A\bar{\gamma}\bar{C}] \rightarrow C^\rho \bar{C}AA\bar{\gamma}\bar{C} & [\bar{C}\bar{\gamma}B] \rightarrow \bar{C}C^\rho \bar{C}A\beta B \end{array}$$

where  $\bar{B}^{r\text{ho}} = (\bar{B}B)^q (A\bar{A})^p \bar{B}$  and  $C^\rho = (C\bar{C})^s (\bar{A}A)^r C$ .

□

### 3.6 The forms $\left| \begin{array}{ccc} 3^+ & 3^+ & 1 \\ 2^+ & 0 & 2^+ \end{array} \right|$ and $\left| \begin{array}{ccc} 2^+ & 2^+ & 0 \\ 2^+ & 0 & 2^+ \end{array} \right|$

In fact there are some lovely geometric methods for more completely understanding these two forms. In the meanwhile, however, through the same methods we have been applying so far, we prove:

#### Theorem 3.6

$$\left[ \begin{array}{ccc} 5^+ & 5^+ & 1 \\ 4^+ & 0 & 6^+ \end{array} \right] \quad \left[ \begin{array}{ccc} 5^+ & 7^+ & 1 \\ 6^+ & 0 & 4^+ \end{array} \right] \quad \left[ \begin{array}{ccc} 4^+ & 8^+ & 0 \\ 4^+ & 0 & 8^+ \end{array} \right] \quad \left[ \begin{array}{ccc} 6^+ & 6^+ & 0 \\ 4^+ & 0 & 8^+ \end{array} \right]$$

It is worth noting that we can construct lifts of our original productions on triangles that will precisely capture all the structure we use here. Indeed, in principle, this is not intrinsically difficult, but would be exceedingly tedious.

Following a few more minor observations, we will have completely sorted out the forms  $\begin{bmatrix} 2^+ & 2^+ & 0 \\ 0 & 0 & 2^+ \end{bmatrix}$  and  $\begin{bmatrix} 3^+ & 3^+ & 1 \\ 1 & 1 & 3^+ \end{bmatrix}$ , leaving a little bit of  $\begin{bmatrix} 3^+ & 3^+ & 1 \\ 0 & 0 & 2^+ \end{bmatrix}$ .

### 3.7.1 $\begin{bmatrix} 2^+ & 2^+ & 0 \\ 0 & 0 & 2^+ \end{bmatrix}$

We begin with a regular production system that models regular  $p$ -gons meeting  $q$ -to-a-vertex, with  $p, q \geq 4$ . We will then take a certain lift of this system that is useful for:

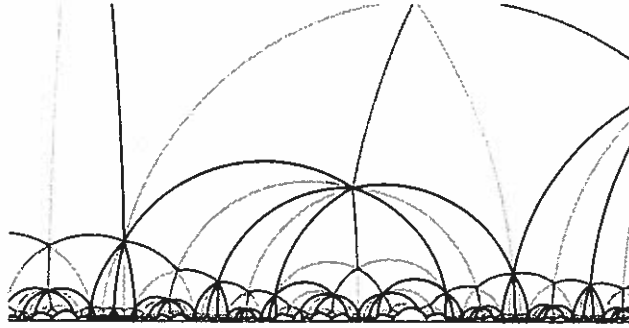
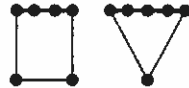


Figure 5: An example of  $\begin{bmatrix} 4^+ & 4^+ & 0 \\ 0 & 0 & 4^+ \end{bmatrix}$

### Theorem 3.7 $\begin{bmatrix} 4^+ & 4^+ & 0 \\ 0 & 0 & 4^+ \end{bmatrix}$

Note that with Lemma 1.1, this completely settles the form  $\begin{bmatrix} 2^+ & 2^+ & 0 \\ 0 & 0 & 2^+ \end{bmatrix}$ .

**Proof** We will prove  $\begin{bmatrix} 2^i & 2^j & 0 \\ 0 & 0 & 2^k \end{bmatrix}$ ,  $i, j, k \geq 2$  by constructing a subdivision of the tiling by  $(2k)$ -gons meeting  $(i+j)$ -to-a-vertex. Let  $p = 2k$ ,  $q = i + j$  and consider the alphabet  $\square, \nabla$  representing the configurations



shown here with  $p = 6$ .

Take the free production system on  $\square, \nabla$  defined by

$$\square \mapsto \nabla^{q-4} \square (\nabla^{q-3} \square)^{p-4}$$

$$\nabla \mapsto \nabla^{q-4} \square (\nabla^{q-3} \square)^{p-3}$$

encoding the layering



shown here with  $p = q = 6$ . Clearly orbits in the system correspond to tilings by abstract  $p$ -gons meeting  $q$ -to-a-vertex.

We now take the following lift, with alphabet  $\square, \square, \nabla, \nabla$  and obvious map onto  $\square, \nabla$ , representing the configurations

So for example, with  $i = j = 3, k = 1$ , the production for  $\nabla$  encodes



As before by inducting on superwords one may easily verify that at every vertex in any complex represented by an orbit in this system, we have a vertex arrangement in  $\mathcal{V}((2i+1)(2j+1)1)$  or  $\mathcal{V}(112k)$ ; and that every vertex in  $\mathcal{V}((2i+1)(2j+1)1)$  meets exactly one edge in  $\mathcal{E}$ . Orbits in the system correspond to  $\mathcal{V}$ -complices with exactly one edge meeting each vertex in  $\mathcal{V}((2i+1)(2j+1)1)$ , and we may apply Lemma 2.1 to obtain  $\begin{bmatrix} 2i+1 & 2j+1 & 1 \\ 1 & 0 & 2k \end{bmatrix}$ .  $\square$

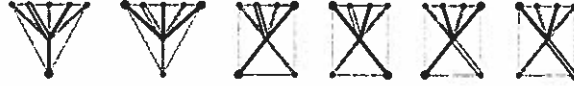
**3.8.1**  $\begin{vmatrix} 3^+ & 3^+ & 1 \\ 1 & 1 & 3^+ \end{vmatrix}$

The cases  $\begin{vmatrix} p & q & 1 \\ 1 & 1 & 3 \end{vmatrix}$ ,  $p, q$  odd, and  $\begin{vmatrix} 5^+ & 3 & 1 \\ 1 & 1 & 5^+ \end{vmatrix}$  are open, for the moment. Other than those, with Corollary 2.2 and Theorem 2.4, the following sorts out the case  $\begin{vmatrix} 3^+ & 3^+ & 1 \\ 1 & 1 & 3^+ \end{vmatrix}$ :

**Theorem 3.10**  $\begin{vmatrix} 5^+ & 5^+ & 1 \\ 1 & 1 & 5^+ \end{vmatrix}$

**Proof** As in Theorem 3.9, we lift the production of Theorem 3.7 in order to mark a set of edges  $\mathcal{E}$  so we may apply Lemma 2.1. We show  $\begin{bmatrix} 2i+1 & 2j+1 & 1 \\ 1 & 1 & 2k+1 \end{bmatrix}$ ,  $i, j, k \geq 2$ . We use the following alphabet:

$\nabla, \nabla, \square, \square, \square,$  and  $\square$  representing the configurations



We use the following productions:

$$\begin{aligned} \nabla &\mapsto \nabla^{j-2} \nabla^{i-2} \square ( \nabla^{j-1} \nabla^{i-2} \square \nabla^{j-2} \nabla^{i-1} \square )^{k-2} \nabla^{j-1} \nabla^{i-2} \square \\ \nabla &\mapsto \nabla^{j-2} \nabla^{i-2} \square ( \nabla^{j-2} \nabla^{i-1} \square \nabla^{j-1} \nabla^{i-2} \square )^{(k-2)} \nabla^{j-2} \nabla^{i-1} \square \\ \square &\mapsto \nabla^{j-2} \nabla^{i-2} \square ( \nabla^{j-2} \nabla^{i-1} \square \nabla^{j-1} \nabla^{i-2} \square )^{(k-2)/2} \\ \square &\mapsto \nabla^{j-2} \nabla^{i-2} \square ( \nabla^{j-1} \nabla^{i-2} \square \nabla^{j-2} \nabla^{i-1} \square )^{(k-2)/2} \\ \square &\mapsto \nabla^{j-2} \nabla^{i-2} \square ( \nabla^{j-2} \nabla^{i-1} \square \nabla^{j-1} \nabla^{i-2} \square )^{(k-2)/2} \\ \square &\mapsto \nabla^{j-2} \nabla^{i-2} \square ( \nabla^{j-1} \nabla^{i-2} \square \nabla^{j-2} \nabla^{i-1} \square )^{(k-2)/2} \end{aligned}$$

Once again, it is not hard to show that in any complex corresponding to orbits in this system, all vertices are in  $\mathcal{V}(\begin{smallmatrix} 2i+1 & 2j+1 & 1 \\ 1 & 1 & 2k+1 \end{smallmatrix})$  and are incident to exactly one edge in  $\mathcal{E}$ . Hence we may apply Lemma 2.1.  $\square$

### 3.9 Remaining triangles

We have not discussed production systems for vertex arrangements of fewer than six triangles; though we did examine many tilings with these arrangements in Section 2 and Section 3.2 of [16]. With some extra care, it is however, perfectly possible to define production systems that capture such vertex arrangements.

We have not explicitly discussed tilings by isocles or equilateral triangles. These are actually much simpler to describe than scalene triangles, and we leave this as an extended exercise.

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